

Quantum Graphs which Optimize the Spectral Gap



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Joint work with Guillaume Lévy, Université Pierre et Marie Curie, Paris
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Eigenvalue optimization problems (for domains)

Fixing the topology, total volume and boundary conditions,
we seek for the shape which maximizes\minimizes an eigenvalue.

Simply connected domains

Faber-Krahn [Dirichlet conditions]: the ball minimizes λ_1 (*no sense maximizing*).

Krahn-Szegö [Dirichlet conditions]: No minimizer for λ_2 ,
but union of two balls serves as an *infimizer*.

Szegö-Weinberger [Neumann conditions]: the ball maximizes λ_1 (*no sense minimizing*).

Multi connected domains

Payne-Weinberger: Planar domains with a single hole,

Dirichlet on outer boundary and Neumann on inner.

Fixing total area and length of outer boundary - annulus (concentric circles) maximizes λ_1 .

More works by: Ashbaugh-Chatelain, Ashbaugh-Benguria, Exner-Mantile, Flucher,
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Outline

Introduction

Infimizers

Supremizers

- Upper bounds

- Spectral gap as a simple eigenvalue

- Gluing graphs

Summary & Conjectures

From a Discrete graph to a Quantum graph

\mathcal{G} a discrete graph with $E < \infty$ edges and $V < \infty$ vertices. Space of edge lengths:

$$\mathcal{L}_{\mathcal{G}} := \left\{ (l_1, \dots, l_E) \in \mathbb{R}^E \mid \sum_{e=1}^E l_e = 1 \text{ and } \forall e, l_e > 0 \right\}$$

$\Gamma(\mathcal{G}; \underline{l})$ denotes the metric graph obtained from \mathcal{G} with edge lengths $\underline{l} \in \mathcal{L}_{\mathcal{G}}$.

Namely, the e^{th} edge corresponds to an interval $[0, l_e]$

Consider the following eigenvalue equation on each $[0, l_e]$: $-\frac{d^2}{dx_e^2} f|_e = k^2 f|_e$,

with the Neumann (Kirchhoff) vertex conditions:

$$\text{Continuity} \quad \forall e_1, e_2 \sim v; \quad f|_{e_1}(v) = f|_{e_2}(v)$$

$$\text{Vanishing sum of derivatives} \quad \sum_{e \sim v} \frac{d}{dx_e} f \Big|_e(v) = 0$$

The spectrum, $\{k_n^2\}_{n=1}^{\infty}$ is discrete and bounded from below:

$$0 = k_0 < k_1 \leq k_2 \leq \dots$$

We call k_1 the **spectral gap** of the graph.

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Spectral gap dependence on edge lengths

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which leads to consider also $\underline{l} \in \partial\mathcal{L}_{\mathcal{G}}$ (some edge lengths vanish),

possibly changing the topology of $\Gamma(\mathcal{G}; \underline{l})$.

Definition 1.

- $\Gamma(\mathcal{G}; \underline{l}^*)$ a maximizer of \mathcal{G} if $\underline{l}^* \in \mathcal{L}_{\mathcal{G}}$ and $k_1[\Gamma(\mathcal{G}; \underline{l}^*)] \geq k_1[\Gamma(\mathcal{G}; \underline{l})]$, $\forall \underline{l} \in \mathcal{L}_{\mathcal{G}}$.
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 - Same definitions for minimizer and infimizer.
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- Supremizer and infimizer always exist.

What about maximizer\minimizer?
 - Which graphs are spectral gap optimizers?



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Spectral gap dependence on edge lengths

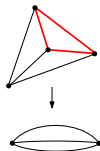
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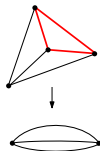
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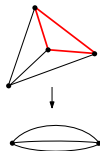
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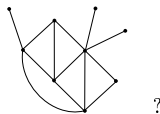
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Quantum Graphs which Optimize the Spectral Gap

- Supremizer and infimizer always exist. What about maximizer\minimizer?
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A few examples

Star graph with $E \geq 2$ edges

Infimum (no minimum): $k_1(1, 0, \dots, 0) = \pi$,

Maximum: $k_1(1/E, \dots, 1/E) = \frac{E}{2}\pi$ (equilateral star)

(Recall: total edge length = 1)

Flower graph with $E \geq 2$ edges

Infimum (no minimum): $k_1(1, 0, \dots, 0) = 2\pi$,

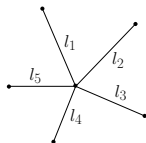
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[Kennedy, Kurasov, Malenová, Mugnolo '16]

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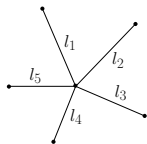
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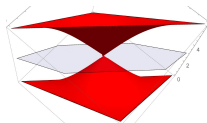
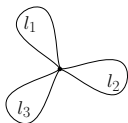


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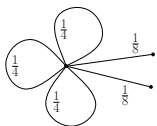
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Quantum Graphs which Optimize the Spectral Gap

A few examples (continued)



Stower (Flétoile) graph with E_p petals, E_l leaves

Infimum (no minimum): $k_1(0 \dots, 0, 1) = \pi$,

Maximum: $k_1(\underline{l}) = (E_p + \frac{E_l}{2})\pi$,

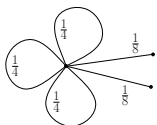
where $\underline{l} = \frac{1}{2E_p + E_l}(\underbrace{2, \dots, 2}_{E_p}, \underbrace{1, \dots, 1}_{E_l})$ (“equilateral” stower),

assuming $E_p + E_l \geq 2$ and $(E_p, E_l) \notin (1, 1)$. [Shown in future slide].

This generalizes stars and flowers results.

Quantum Graphs which Optimize the Spectral Gap

A few examples (continued)



Stower (Flétoile) graph with E_p petals, E_l leaves

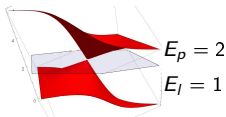
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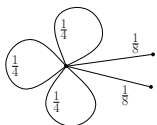


Infimum: $k_1(0, 0, 1) = \pi$,

Maximum: $k_1(\frac{2}{5}, \frac{2}{5}, \frac{1}{5}) = 2\frac{1}{2}\pi$

Quantum Graphs which Optimize the Spectral Gap

A few examples (continued)



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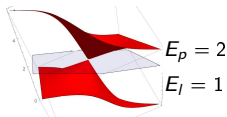
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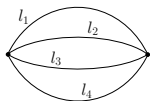


Continuous family of infima: $k_1(0, t, 1 - t) = \pi$,

Continuous family of maxima: $k_1(1 - 2t, t, t) = 2\pi$

Quantum Graphs which Optimize the Spectral Gap

A few examples (continued)



$k_1(l_1, l_2, l_3)$

Mandarin graph with E edges

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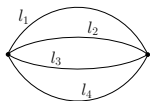
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Length dependence figures - courtesy of Lior Alon

- Which graphs have not only supremizer\infimizer, but also maximizer\minimizer?
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Lower bounds - Known results

$$k_1[\Gamma] \geq \pi$$

with equality iff Γ is a single edge [Nicaise '87; Friedlander '05; Kurasov, Naboko '14].

If Γ has all vertex degrees even then

$$k_1[\Gamma] \geq 2\pi, \quad [\text{Kurasov, Naboko '14}]$$

with a single loop achieving equality (for example).

Remaining questions:

- What about other topologies?
- What are all possible minimizers\infimizers?

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Infimizers - Solution

A **bridge** is an edge whose removal dissects the graph.

Theorem 2 (Band, Lévy).

1. *Let \mathcal{G} be a graph with a bridge. Then*
 - 1.1 *The infimal spectral gap of \mathcal{G} equals π .*
 - 1.2 *The unique infimizer is the unit interval.*
2. *Let \mathcal{G} be a bridgeless graph. Then*
 - 2.1 *The infimal spectral gap of \mathcal{G} equals 2π .*
 - 2.2 *Any infimizer is a symmetric necklace graph.*

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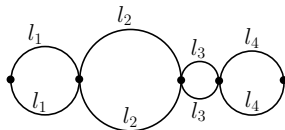


Figure: symmetric necklace graph

Infimizers - Solution

A **bridge** is an edge whose removal disconnects the graph.

Theorem 2 (Band, Lévy).

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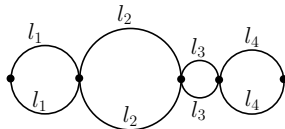


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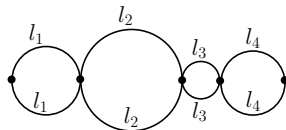


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Upper bounds - Known results

- Global bound

$$k_1[\Gamma] \leq E\pi,$$

equality if and only if Γ is an equilateral mandarin or equilateral flower [Kennedy, Kurasov, Malenová, Mugnolo '16].

This fully answers optimization for flowers and mandarins:

supremizers (also maximizers) are equilateral.

- If Γ is a tree then

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Upper bounds - Further progress

Proposition 3 (Band, Lévy).

If Γ is a tree with E_l leaves then $k_1[\Gamma] \leq \frac{E_l}{2}\pi$.

Proof idea.

$d(\Gamma) := \max\{d(x, y) \mid x, y \in \Gamma\}$ graph diameter.

Combine $k_1[\Gamma] \leq \frac{\pi}{d(\Gamma)}$ with $d(\Gamma) \geq \frac{2}{E_l}$ (the latter true for trees).

□

Proposition 4 (Band, Lévy).

Let \mathcal{G} be a graph with E edges, out of which E_l are leaves.

If $(E, E_l) \notin \{(1, 1), (1, 0), (2, 1)\}$ then $\forall \underline{l} \in \mathcal{L}_{\mathcal{G}}, \quad k_1[\Gamma(\mathcal{G}; \underline{l})] \leq \pi \left(E - \frac{E_l}{2}\right)$.

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Take Γ and attach two vertices to obtain Γ' (illegal move in our game). Get $k_1(\Gamma) \leq k_1(\Gamma')$.

Repeatedly attach all inner vertices to obtain a stower with E_l leaves and $E - E_l$ petals.

Use bound on stowers: $k_1[\Gamma] \leq \pi \left(E - \frac{E_l}{2}\right)$ [to appear in a future slide]



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Spectral gap as a simple eigenvalue - Critical points

Try to find supremizers by seeking for local critical points in $\mathcal{L}_{\mathcal{G}}$.

Derivatives with respect to edge lengths may be calculated for simple eigenvalues.

Theorem 5 (Band, Lévy).

Let \mathcal{G} be a discrete graph and $\underline{l} \in \mathcal{L}_{\mathcal{G}}$.

Assume that $\Gamma(\mathcal{G}; \underline{l})$ is a supremizer of \mathcal{G} with simple spectral gap $k_1[\Gamma(\mathcal{G}; \underline{l})]$.

Then $\Gamma(\mathcal{G}; \underline{l})$ is not a unique supremizer:

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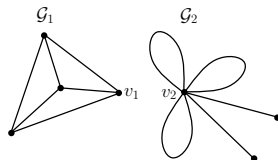
Gluing graphs - Vertex connectivity one

Let $\mathcal{G}_1, \mathcal{G}_2$ be discrete graphs, and v_i ($i = 1, 2$) be a vertex of \mathcal{G}_i .

Let \mathcal{G} be the graph obtained by identifying (gluing) v_1 and v_2 .

If we know the supremizers Γ_1, Γ_2 of $\mathcal{G}_1, \mathcal{G}_2$,

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Let $\mathcal{G}_1, \mathcal{G}_2$ be discrete graphs.

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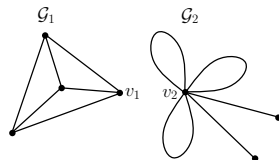
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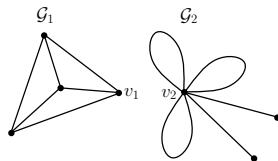
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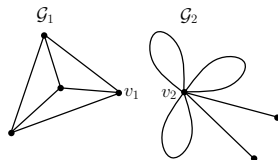
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Gluing graphs - Corollaries

Corollary 7.

Let \mathcal{G} be a stower with $E_p + E_l \geq 2$ and $(E_p, E_l) \neq (1, 1)$. Then a maximizer is the “equilateral” stower graph with spectral gap $\pi \left(E_p + \frac{E_l}{2} \right)$.

This maximizer is unique for $(E_p, E_l) \notin \{(2, 0), (1, 2)\}$.

Proof idea.

Prove the statement for “small” stowers. Then glue them to construct any stower. □

Recall

Proposition 4:

Let \mathcal{G} be a graph with E edges, out of which E_l are leaves.

If $(E, E_l) \notin \{(1, 1), (1, 0), (2, 1)\}$ then $\forall \underline{l} \in \mathcal{L}_{\mathcal{G}}, \quad k_1[\Gamma(\mathcal{G}; \underline{l})] \leq \pi \left(E - \frac{E_l}{2} \right)$.

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- Optimization problem fully solved for infimizers\minimizers.
- Supremizers
 - ▶ Improved upper bounds by conditioning on number of leaves,
 $k_1 \leq \pi \left(E - \frac{E_1}{2} \right)$ (global) and $k_1 \leq \pi \frac{E_1}{2}$ (for trees).
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Supremizer candidates are stowers and mandarins (are there any others?)

\Rightarrow lower bounds on supremal spectral gap

Getting to a stower gives $\pi \left(\beta + \frac{E_l}{2} \right)$,

where $\beta := E - V + 1$ is the graph's first Betti number.

Getting to a mandarin:

Partition vertices $V = V_1 \cup V_2$.

$E(V_1, V_2) := \#$ of edges connecting V_1 to V_2 .

Maximal spectral gap *among all mandarins* is

$\pi \cdot \max_{V_1, V_2} E(V_1, V_2)$. (Cheeger-like constant)

Compare $\pi \left(\beta + \frac{E_l}{2} \right)$ (stower) with $\pi \cdot \max_{V_1, V_2} E(V_1, V_2)$ (mandarin).

$E(V_1, V_2) = \beta + 1 - (\beta_1 + \beta_2)$, where β_i is the Betti number of V_i graph.

If $E_l \leq 1$ then mandarin wins if and only if we find $\beta_1 = \beta_2 = 0$.

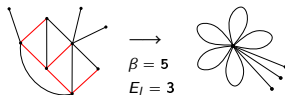
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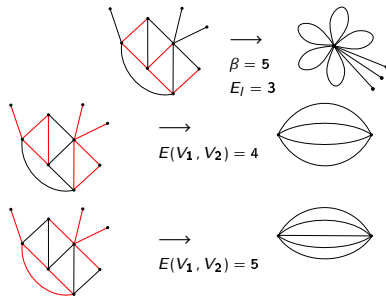
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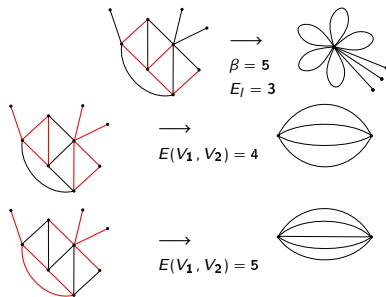
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Maximal spectral gap *among all mandarins is*

$\pi \cdot \max_{V_1, V_2} E(V_1, V_2)$. (Cheeger-like constant)

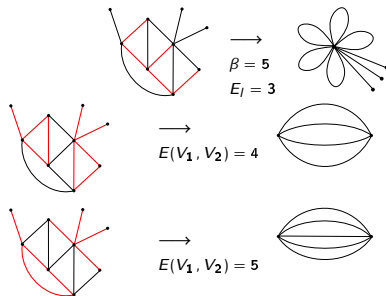
Compare $\pi \left(\beta + \frac{E_l}{2} \right)$ (stower) with $\pi \cdot \max_{V_1, V_2} E(V_1, V_2)$ (mandarin).

$E(V_1, V_2) = \beta + 1 - (\beta_1 + \beta_2)$, where β_i is the Betti number of V_i graph.

If $E_l \leq 1$ then mandarin wins if and only if we find $\beta_1 = \beta_2 = 0$.

If $E_l \geq 2$ then mandarin never wins (possibility for a tie).

Leads to conjectures....



Conjectures

- Supremizer is either a mandarin or a stower.
- Supremum is obtained when order of symmetry group is maximized.
- Supremum is obtained when multiplicity of spectral gap is maximized.

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Quantum Graphs which Optimize the Spectral Gap



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(arXiv:1608.00520)

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